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We have overwritten the address of the memory: it could have been anywhere, we don't know anymore

That area of memory is now *garbage*. It takes up space but the program can't get at it

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The important thing is to ensure a pointer to every allocated chunk is somehow accessible (directly or indirectly) by the program and can be accessed or freed if necessary

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Memory leaks often go unnoticed as programmers often test their programs on small examples: small enough that the amount of garbage is still small and malloc always succeeds

They only discover the error when their code goes into production on big examples and then starts failing

Aside. Current operating systems clean up after you when your program exits, returning all malloced memory. Some early operating systems didn't, meaning poorly written programs could jam up the entire computer, eventually requiring a reboot

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Tools like valgrind will tell you how much memory you have malloced and not freed

malloc and free are a major source of bugs in C programs

Using memory you have not malloced

- Using memory you have not malloced
- freeing memory more than once

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- And so on

On the other hand, malloc and free are extremely useful in the right hands

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- they concentrate the programmer's attention towards the efficient use of memory
- they are reasonably fast
- the programmer can tune their use to the problem in hand

Exercise. What is the bug here?

```
int a[10];
...
free(a);
```

Exercise. malloc and free are fast, but not free: they take some time (and some overhead space) to manage memory. Find out how much of an overhead they incur on your computer

Exercise. Compare this with Java's memory management

Exercise. Look up alloca and dynamic stack allocation

Exercise. Deliberately write bad code that does these kinds of things. Run it and see what goes wrong. Use valgrind on your code

Exercise. Deliberately write good code that avoids these kinds of things

Exercise. Think about the symmetry:

```
int *a = malloc(...);
free(a);
```

giving a pointer that points at a non-object; and

```
int *a = malloc(...);
a = malloc(...);
```

giving an object that no pointer pointing to it

Having malloc and free is simultaneously one of the great strengths of C and one of its great weaknesses

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Java programs generate garbage at a prodigious rate

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So Java includes (as part of the Java system) a *garbage* collector that periodically trawls through memory looking for inaccessible garbage: chunks of memory that can never be accessed in the program as the program has overwritten/lost the pointers to those chunks

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It collects the areas of garbage memory together and then can allocate those bytes in subsequent calls

Note it is safe to reallocate those bytes as by definition garbage is inaccessible to the program, thus reusing them can have no effect on the program

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Automatic memory management:

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It a choice of trade-offs

- Releases the programmer from having to worry about memory
- Should never go wrong
- Encourages sloppy programming
- Has a significant time and space overhead in management and garbage collection

Manual memory management:

Requires the programmer to think carefully about memory usage

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- Encourages careful use of memory
- Can be tuned for a specific application
- Is a frequent source of errors

It's your choice and should be taken into account when you are choosing a programming language to implement a project

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C does have bolt-on garbage collectors, if you really want them

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The reservation exists until we do a free

This means this kind of code is OK (and very common):

```
struct intlist *make(int v)
{
   struct intlist *newl;
   ...
   newl = (struct intlist *)malloc(sizeof(struct intlist));
   ...
   return newl;
}
```

as the reservation persists beyond the end of the function call make, so the returned pointer remains valid outside of the function call

However, this is bad:

```
struct intlist *make(int v)
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By "exist" we mean "is valid". It's still there in memory!

Because the pointer returned is still a pointer to somewhere in memory, the code might even work for a while

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Until you have another function call that extends the stack again to cover that place where your structure lives. And then overwrites it with whatever locals and arguments that function requires

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More precisely: it is OK to use pointers to things on the stack while that frame is still active. Thus using such a pointer within the current function is fine; as is passing the pointer down to "deeper" functions. But you must never return the pointer up to a place where the frame has gone

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Exercise. Investigate to see what happens when you return pointers to things on the stack

Exercise. What about

```
struct intlist make(int v)
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   return newl;
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You will need to manipulate files: read and write data

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C provides two principal kinds of access to files:

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We shall look at buffered I/O: it's the one you will use the most

If you need unbuffered I/O, you will easily be able to pick it up for yourself

The major operations on files are

- open, close (functions fopen, fclose)
- read, write (functions fread, fwrite, fprintf)

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In fact, you will always be using a pointer to a FILE, a FILE*

```
int main(int argc, char *argv[])
₹
  FILE *in, *out;
  char buf[1024]; // chunk of bytes
  int nread;
  in = fopen(argv[1], "r"); // ought to check
  out = fopen(argv[2], "w"); // for success
  do ₹
   nread = fread(buf, 1, 1024, in);
    fwrite(buf, 1, nread, out);
  } while (nread > 0);
  fclose(in);
  fclose(out);
  return 0;
```

#include <stdio.h>

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- We should check for success of both fopens. They will return NULL if they failed. For example, trying to read a file that does not exist or we do not have permission to read

```
nread = fread(buf, 1, 1024, in);
fwrite(buf, 1, nread, out);
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 We repeatedly read bytes from in. We shall try to read 1024 items of size 1 byte each into the buffer buf

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- We write that number of bytes to out
- We repeat until there are no more bytes to read

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- We then close in and out
- It is important to close files, particularly when writing, to ensure all the data is safely written to disk before the program ends
- Also, fclose does a free of the relevant datastructures that fopen made behind the scenes

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Using command-line shells we can redirect the two kinds of output to different places

```
fwrite(str, 1, 12, stdout);
is an unlikely way of writing a string to the screen
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Exercise. Look at the man pages for these file functions, particularly fopen

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In fact, printf is the same as fprintf(stdout, ...)

And fprintf(stderr, ...) is the way you usually report errors to the user

Exercise. Look at fscanf (and scanf), the "opposite" to printf that reads text formatted input

Exercise. Make sure you understand the distinction between using fread to read a (4 byte, say) integer and using fscanf to read a (character string) integer

Exercise. Look up feof, fflush and ferror

Exercise. Read man stdio